

Entropy based data compression algorithms

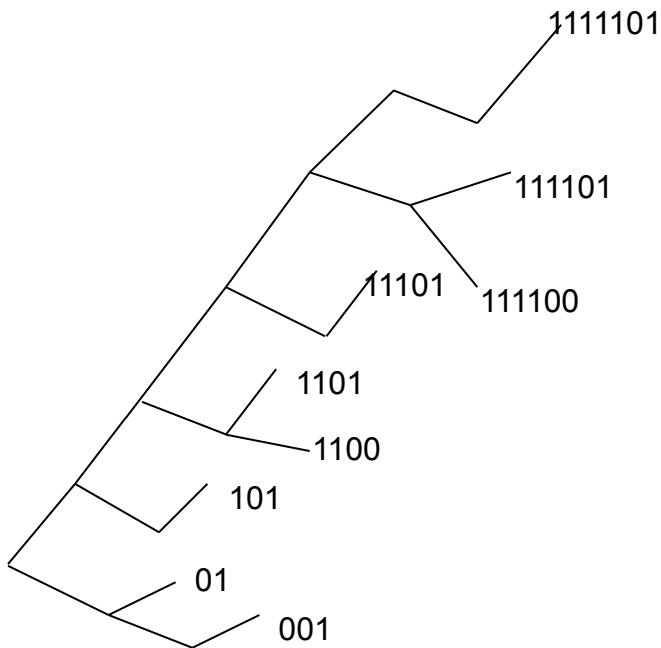
Shannon-Fano, Huffman,
Shannon-Fano-Elias codes

Shannon – Fano code

$$p_1 = 0.49; p_2 = 0.14; p_3 = 0.14; p_4 = 0.07; p_5 = 0.07; p_6 = 0.04; p_7 = 0.02; p_8 = 0.02; p_9 = 0.01$$

$$l_1 = \left\lceil \log_2 \frac{1}{p_1} \right\rceil = \lceil 1.029 \rceil = 2; \quad l_2 = \left\lceil \log_2 \frac{1}{p_2} \right\rceil = \lceil 2.836 \rceil = 3; \quad l_3 = \left\lceil \log_2 \frac{1}{p_3} \right\rceil = \lceil 2.836 \rceil = 3; \quad l_4 = \left\lceil \log_2 \frac{1}{p_4} \right\rceil = \lceil 3.836 \rceil = 4; \quad l_5 = \left\lceil \log_2 \frac{1}{p_5} \right\rceil = \lceil 3.836 \rceil = 4;$$

$$l_6 = \left\lceil \log_2 \frac{1}{p_6} \right\rceil = \lceil 4.64 \rceil = 5; \quad l_7 = \left\lceil \log_2 \frac{1}{p_7} \right\rceil = \lceil 5.64 \rceil = 6; \quad l_8 = \left\lceil \log_2 \frac{1}{p_8} \right\rceil = \lceil 5.64 \rceil = 6; \quad l_9 = \left\lceil \log_2 \frac{1}{p_9} \right\rceil = \lceil 6.64 \rceil = 7$$



Szimb.	Kódszó
X_1	01
X_2	001
X_3	101
X_4	1100
X_5	1101
X_6	11101
X_7	111100
X_8	111101
X_9	1111101

Efficiency

$$p_1 = 0.49; p_2 = 0.14; p_3 = 0.14; p_4 = 0.07; p_5 = 0.07; p_6 = 0.04; p_7 = 0.02; p_8 = 0.02; p_9 = 0.01$$

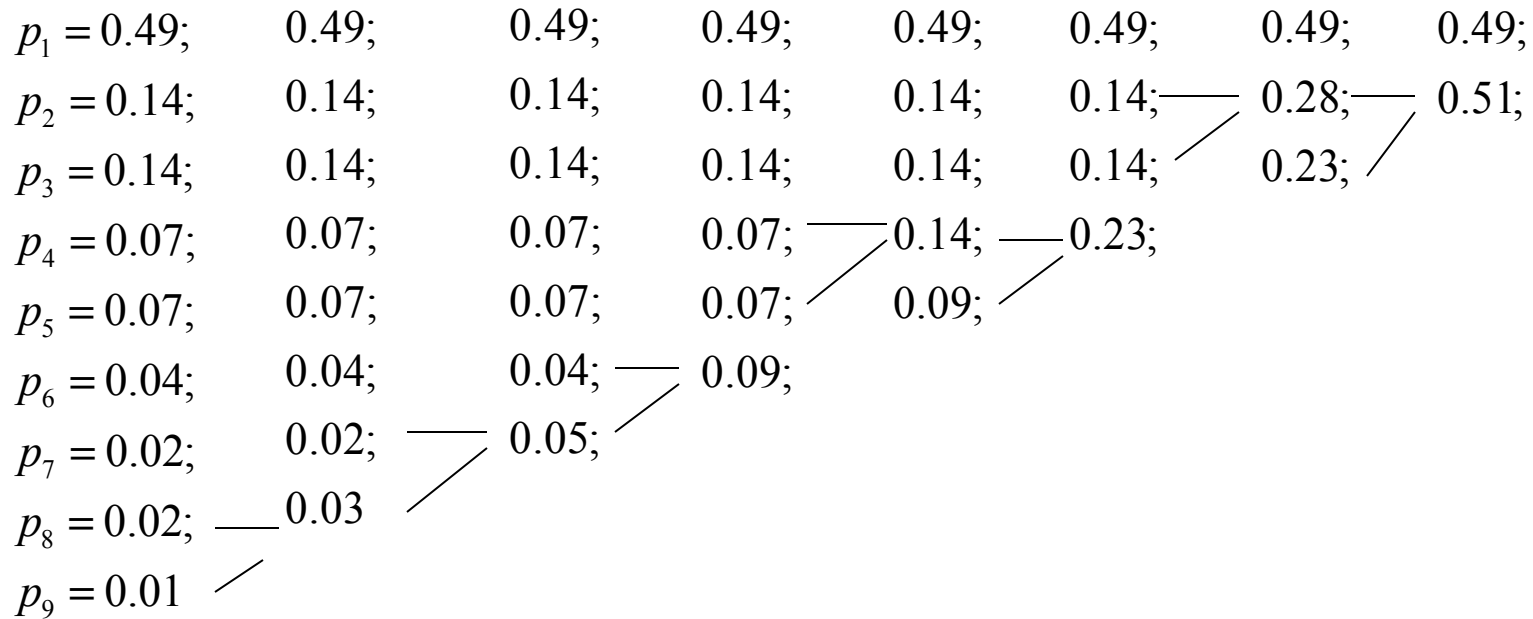
$$l_1 = \left\lceil \log_2 \frac{1}{p_1} \right\rceil = \lceil 1.029 \rceil = 2; \quad l_2 = \left\lceil \log_2 \frac{1}{p_2} \right\rceil = \lceil 2.836 \rceil = 3; \quad l_3 = \left\lceil \log_2 \frac{1}{p_3} \right\rceil = \lceil 2.836 \rceil = 3; \quad l_4 = \left\lceil \log_2 \frac{1}{p_4} \right\rceil = \lceil 3.836 \rceil = 4; \quad l_5 = \left\lceil \log_2 \frac{1}{p_5} \right\rceil = \lceil 3.836 \rceil = 4;$$
$$l_6 = \left\lceil \log_2 \frac{1}{p_6} \right\rceil = \lceil 4.64 \rceil = 5; \quad l_7 = \left\lceil \log_2 \frac{1}{p_7} \right\rceil = \lceil 5.64 \rceil = 6; \quad l_8 = \left\lceil \log_2 \frac{1}{p_8} \right\rceil = \lceil 5.64 \rceil = 6; \quad l_9 = \left\lceil \log_2 \frac{1}{p_9} \right\rceil = \lceil 6.64 \rceil = 7$$

$$L = 0.49 * 2 + 0.28 * 3 + 0.14 * 4 + 0.04 * 5 + 0.04 * 6 + 0.01 * 7 =$$
$$= 0.98 + 0.84 + 0.56 + 0.2 + 0.24 + 0.07 = 2.89$$

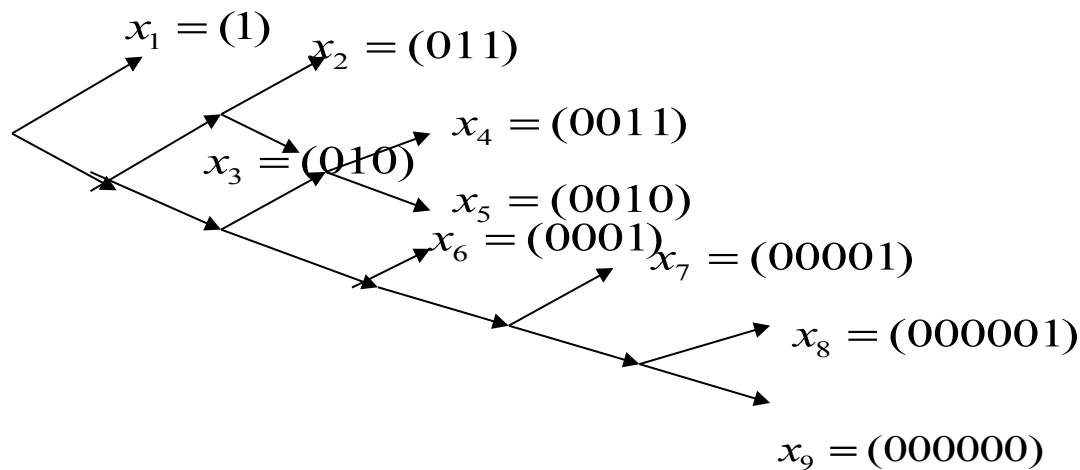
$$H(X) = \sum_{i=1}^9 p_i \log_2 \left(\frac{1}{p_i} \right) = 2.314$$

$$\text{Efficiency} = \frac{H(X)}{L} = 0.8 = 80\%$$

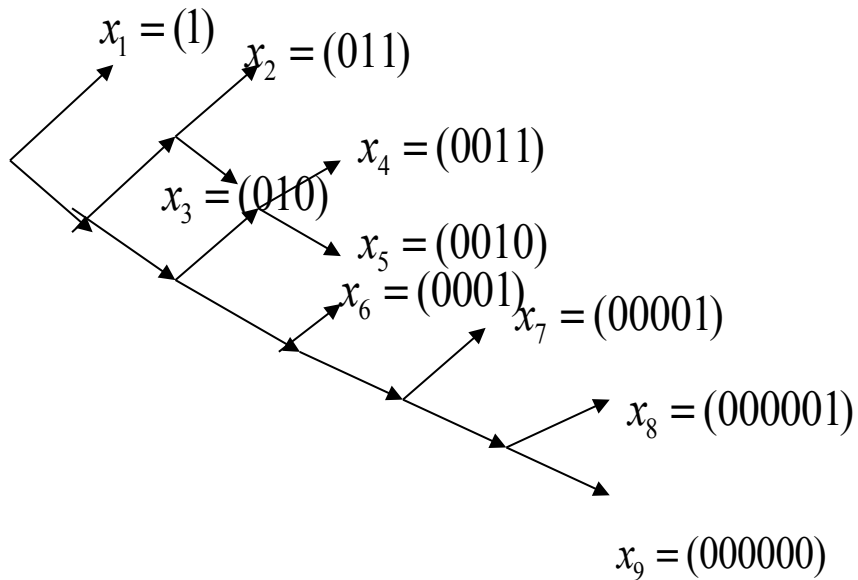
Solution by Huffman coding



$p_1 = 0.49;$	0.49;	0.49;	0.49;	0.49;	0.49;	0.49;	0.49;	0.49;	1
$p_2 = 0.14;$	0.14;	0.14;	0.14;	0.14;	0.14;	0.28;	0.51;		
$p_3 = 0.14;$	0.14;	0.14;	0.14;	0.14;	0.14;	0.23;			
$p_4 = 0.07;$	0.07;	0.07;	0.07;	0.14;	0.23;				
$p_5 = 0.07;$	0.07;	0.07;	0.07;	0.09;					
$p_6 = 0.04;$	0.04;	0.04;	0.09;						
$p_7 = 0.02;$	0.02;	0.05;							
$p_8 = 0.02;$	0.03								
$p_9 = 0.01$									



The coding LUT



Szimb.	Kódszó
X_1	1
X_2	011
X_3	010
X_4	0010
X_5	0011
X_6	0001
X_7	00001
X_8	000001
X_9	000000

Efficiency

$$L^{HUFF} = 0.49 * 1 + 0.14 * 3 + 0.14 * 3 + 0.07 * 4 + 0.07 * 4 + 0.04 * 4 + 0.02 * 5 + 0.02 * 6 + 0.01 * 6 = 2.33$$

$$L^{SF} = 0.49 * 2 + 0.28 * 3 + 0.14 * 4 + 0.04 * 5 + 0.04 * 6 + 0.01 * 7 = 2.89$$

$$f_s = 160 \text{ MHz} = 1,6 * 10^8 \quad \text{esetén}$$

$$R_{Huffman} = 372,8 \text{ Mbps}$$

$$R_{SF} = 462 \text{ Mbps}$$

$$\text{Diff. dataspeed} = 1,6 * 10^8 (2.89 - 2.33) = 89,6 * 10^6 \approx 90 \text{ Mbps}$$

Solution by Shannon-Fano-Elias codes

Prob.	Mass func. Distribution	Mod. Distr.	Bin. conversion	Codeword
$p_1 = 0.49;$	$F(1) = 0;$	$\bar{F}(1) = 0.245;$	$0.00111110101110000101000;$	001;
$p_2 = 0.14;$	$F(2) = 0.49;$	$\bar{F}(2) = 0.56;$	$0.1000111101101110000101;$	1000;
$p_3 = 0.14;$	$F(3) = 0.63;$	$\bar{F}(3) = 0.7;$	$0.101\overline{100};$	1011;
$p_4 = 0.07;$	$F(4) = 0.77;$	$\bar{F}(4) = 0.805;$	$0.11001110000101000111101;$	11001;
$p_5 = 0.07;$	$F(5) = 0.84;$	$\bar{F}(5) = 0.875;$	$0.11110;$	11100;
$p_6 = 0.04;$	$F(6) = 0.91;$	$\bar{F}(6) = 0.93;$	$0.11101110000101000111101;$	111011;
$p_7 = 0.02;$	$F(7) = 0.95;$	$\bar{F}(7) = 0.96;$	$0.11110101110000101000;$	1111010;
$p_8 = 0.02;$	$F(8) = 0.97;$	$\bar{F}(8) = 0.98;$	$0.111110101110000101000;$	1111101;
$p_9 = 0.01$	$F(9) = 0.99;$	$\bar{F}(9) = 0.995;$	$0.11111110101110000101000;$	11111110;

SFE code (cont')

Szimb.	Kódszó
X_1	001
X_2	1000
X_3	1011
X_4	11001
X_5	11100
X_6	111011
X_7	1111010
X_8	1111101
X_9	11111110

$$L^{SFE} = 0.49*3 + 0.14*4 + 0.14*4 + 0.07*5 + 0.07*5 + 0.04*6 + 0.02*7 + 0.02*7 + 0.01*8 = 3.89$$

$$L^{SF} = 2.89 \quad L^{HUFF} = 2.33$$

Data speed

$$L^{HUFF} = 2.33$$

$$L^{SF} = 2.89$$

$$L^{SFE} = 3.89$$



$$R_{Huffman} = 372,8 \text{ Mbps}$$

$$R_{SF} = 462 \text{ Mbps}$$

$$R_{SFE} = 622,4 \text{ Mbps}$$

Without data compression 640 Mbps !!!!

Conclusion: small improvement in „L” matters a lot in data speed !!!

Comparative analysis

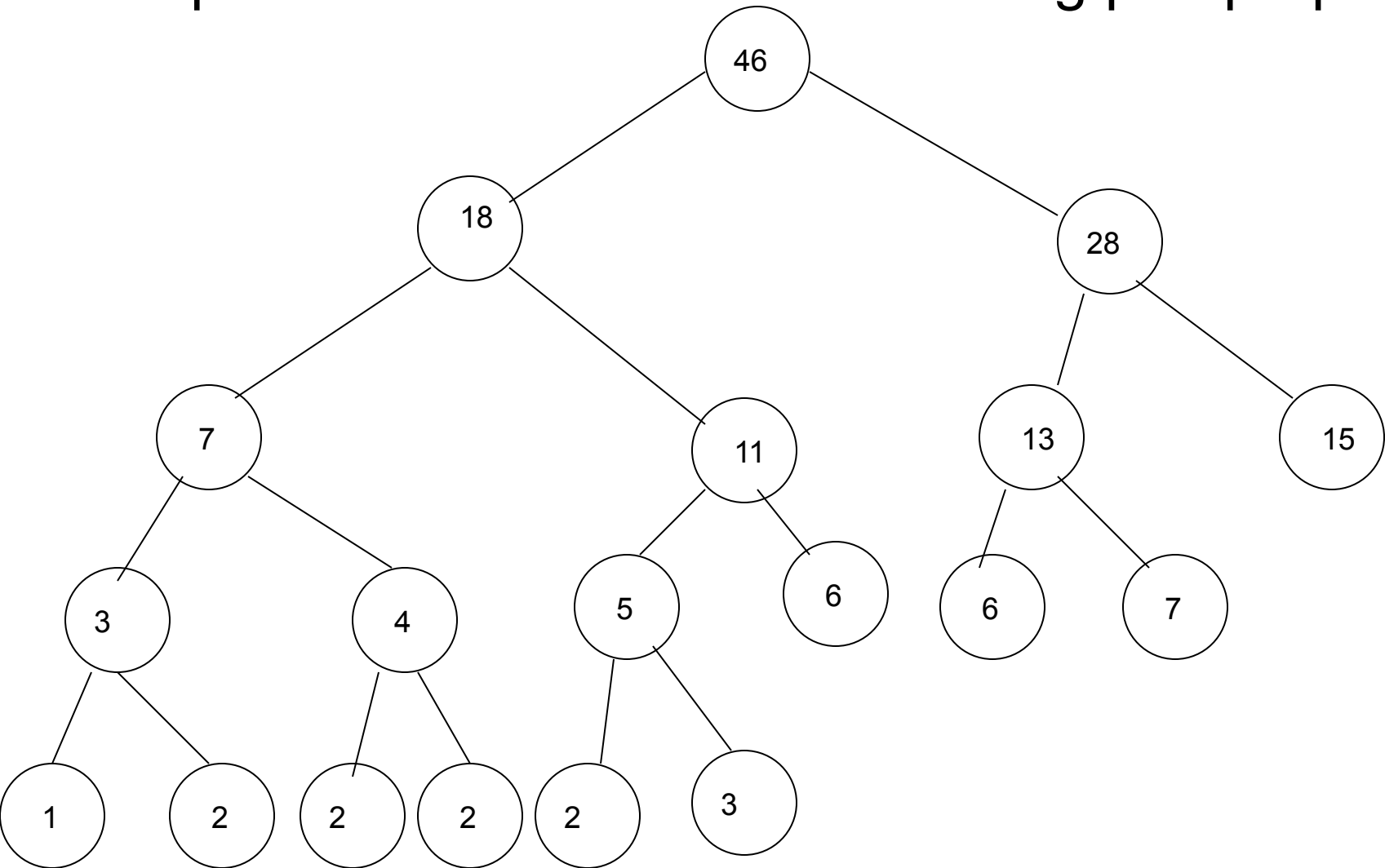
$$f_s = 160 \text{ MHz} = 1,6 * 10^8$$

performance

<i>Kód</i>	<i>Telj. kép</i>	<i>Átl. hossz</i>	<i>Adatátvseb</i>	<i>Alg. Kompl.</i>
Huffman	L_{opt}	2.33	372,8 Mbps	Ker.+ bin. fa
SF	$H(X) < L < H(X) + 1$	2.89	462,4 Mbps	bin. fa
SFE	$H(X) < L < H(X) + 2$	3.89	622,4 Mbps	Bin. konv.

alg. simplicity

Adaptive Huffman codes – the sibling pair property



1,2,2,2,2,3; 3,4,5,6,6,7; 7,11,13,15; 18,28; 46 monotone increasing sequence from bottom layer to top layer

Every parent node is the sum of children nodes

Construction of adaptive Huffman codes

DCDA

