# Dynamic data structures – Linked lists Basics of Programming 1



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Dynamic data structures
Self-referencing structure
Singly linked lists
Definition

- Traversing
- Stack
- Insertion
- Deleting

# Chapter 1

Dynamic data structures



- We are writing a chess program, in which there is undo option for arbitrary number of moves.
- The undo-list is the log of the game, its elements are the moves.
  - Which piece
  - From where
  - Where to
  - Who is captured (removed)
- For logging we use the memory we really need, no more.



- We are writing a chess program, in which there is undo option for arbitrary number of moves.
- The undo-list is the log of the game, its elements are the moves.
  - Which piece
  - From where
  - Where to
  - Who is captured (removed)
- For logging we use the memory we really need, no more.
- The final length of the log will be known only at the end of the game.
- We have to increase the amount of allocated memory with each step (or reduce it, if we undo a move).















If we use realloc for resizing an array, it may cause many unnecessary copying of data.



We need a data structure that does not use continuous blocks. of memory, and its strucure changes dynamically during the lifecycle of the program.

# Dynamic data structure

Dynamic data structure:



- its size or structure changes during the lifecycle of the program
- it is realized with self-referencing structure

### Self-referencing structure

A compound data structure, that contains pointers pointing to itself

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### Self-referencing structure

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next points to a structure that is of the same type, as the one containing the pointer itself.

# Dynamic data structure

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### Self-referencing structure

A compound data structure, that contains pointers pointing to itself

- next points to a structure that is of the same type, as the one containing the pointer itself.
- struct listelem structure is renamed to listelem, but when declaring next, we must use the long name (because the compiler doesn't know, what nickname we will give to it).

## Chapter 2

Singly linked lists



- List of listelem type variables
- Memory is allocated dynamically, separately for each element
- Elements do not form a continuous block in memory
- Each element contains the address of the next element
- The first element is defined by the head pointer
- The last element points to nowhere (NULL)





Empty list

head NULL

# Linked list

Empty list



 List is a self-referencing (recursive) data structure. Each element points to a list.



### List or array



#### The array

- occupies as much memory, as needed for storing the data
- needs a continuous block of memory
- any element can be accessed directly (immediately), by indexing
- inserting a new data involves a lot of copying

### List or array



#### The array

- occupies as much memory, as needed for storing the data
- needs a continuous block of memory
- any element can be accessed directly (immediately), by indexing
- inserting a new data involves a lot of copying
- The list
  - elements store the address of the next element, this may need a lot of memory
  - can make use of gaps in the fragmented memory
  - only the next element can be accessed immediately
  - inserting a new element involves only a little work




























































































 For traversing we need an auxiliary pointer (p), that will run along the list.





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Dynamic Singly linked

# Passing a list to a function



link

As a list is determined by its starting address, we only need to pass the startig address for the function

```
void traverse(listelem *head) {
1
     listelem *p = head;
2
     while (p != NULL)
3
     {
4
       printf("%d ", p->data);
5
6
       p = p - > next;
     }
7
  }
8
```

Dynamic Singly linked

# Passing a list to a function



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     while (p != NULL)
3
     ł
4
       printf("%d ", p->data);
5
       p = p - > next;
6
     }
7
  }
8
```

the same with for loop

```
void traverse(listelem *head) {
    listelem *p;
    for (p = head; p != NULL; p = p->next)
        printf("%d ", p->data);
    }
```

```
p = (listelem*)malloc(sizeof(listelem));
```

- p 2 = 5;
- 3 p->next = head;

```
4 head = p;
```







- p = (listelem\*)malloc(sizeof(listelem));
- 2 p->data = 5;
- 3 p->next = head;
- 4 head = p;





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- $_{2}$  p->data = 5;
- 3 p->next = head;

```
4 head = p;
```



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- $_{2}$  p->data = 5;
- 3 p->next = head;

```
4 head = p;
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- p 2 data = 5;
- 3 p->next = head;

```
4 head = p;
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- 2 p->data = 5;
- 3 p->next = head;

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- p > data = 5;
- 3 p->next = head;

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3 p->next = head;

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```



```
p = (listelem*)malloc(sizeof(listelem));
```

- 2 p->data = 5;
- 3 p->next = head;

```
4 head = p;
```



```
Dynamic Singly linked Def Traversing Stack Insertion Deleting
```

Inserting element to the front of the list, with a fur a presence of the list, with a fur a presence of the list o

 As the starting address is changed when inserting, we have to return it (pass it back)

```
listelem *push_front(listelem *head, int d)
  ł
2
     listelem *p = (listelem*)malloc(sizeof(listelem));
3
     p - > data = d;
4
    p - > next = head;
5
    head = p;
6
     return head;
7
8
  }
                                                             link
```

```
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3
    p - > data = d;
4
  p->next = head;
5
   head = p;
6
    return head;
7
  }
8
                                                         link

    Usage of function

  listelem *head = NULL; /* empty list */
1
  head = push_front(head, 2); /* head is changed! */
2
```

```
3 head = push_front(head, 4);
```



Another option is to pass the starting address by its address

```
void push_front(listelem **head, int d)
{
listelem *p = (listelem*)malloc(sizeof(listelem));
p->data = d;
p->next = *head;
*head = p; /* *head is changes, this is not lost */
}
```

Def Traversing Stack Insertion Deleting



Another option is to pass the starting address by its address

```
void push_front(listelem **head, int d)
{
listelem *p = (listelem*)malloc(sizeof(listelem));
p->data = d;
p->next = *head;
*head = p; /* *head is changes, this is not lost */
}
```

In this case the usage of the function is:

```
1 listelem *head = NULL; /* empty list */
2 push_front(&head, 2); /* calling with address */
3 push_front(&head, 4);
```



- 1 p = head; 2 head = head->next;
- 2 nead = nead->next
- 3 free(p);







- 2 head = head->next;
- 3 free(p);







- 2 head = head->next;
- 3 free(p);



Dynamic Singly linked

Def Traversing Stack Insertion Deleting



- 1 p = head; 2 head = head->next; 3 free(p);
  - 0x1000 0x1234 0x2345 0x3456 1 2 3 4 0x1234 0x2345 0x3456 NULL head 0x1000 р 0x1000

Dynamic Singly linked

Def Traversing Stack Insertion Deleting





- p = head;
  head = head->next;
  frac(r);
- 3 free(p);





- 1 p = head; 2 head = head->next;
- 3 free(p);



2

3







- 1 p = head; 2 head = head->next;
- 3 free(p);



link



```
listelem *pop_front(listelem *head)
1
   ł
2
     if (head != NULL) /* not empty */
3
     {
4
        listelem *p = head;
5
        head = head->next;
6
        free(p);
7
     }
8
     return head;
9
   }
10
```

An empty list must be handled separately

link



```
listelem *pop_front(listelem *head)
1
   ł
2
     if (head != NULL) /* not empty */
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     Ł
4
        listelem *p = head;
5
       head = head->next;
6
       free(p);
7
     }
8
     return head;
9
10
   }
```

- An empty list must be handled separately
- Of course we could use the solution when calling the function with the address of head




What we have so far is already enough for storing the undo-list

```
listelem *head = NULL:
                                  /* empty list */
1
  head = push_front(head, 2); /* step */
2
  head = push_front(head, 4); /* step */
3
  printf("The last inserted element: %d\n", head->data);
4
  head = pop_front(head);
                                /* undo */
5
  head = push_front(head, 5);
                                  /* step */
6
                                  /* step */
  head = pop_front(head);
7
  head = pop_front(head);
                                   /* strep */
8
```

- The stack is a LIFO: Last In, First Out
- We can access the last inserted element first













```
1 for (p = head; p->next != NULL; p = p->next);
2 p->next = (listelem*)malloc(sizeof(listelem));
3 p->next->data = 5;
4 p >next = NULL;
```

```
4 p->next->next = NULL;
```



;

$$_3$$
 p->next->data = 5;



```
1 for (p = head; p->next != NULL; p = p->next);
2 p->next = (listelem*)malloc(sizeof(listelem));
3 p->next->data = 5;
4 p >next = NULL;
```

```
4 p->next->next = NULL;
```





;

$$_{3}$$
 p->next->data = 5;







;

$$_{3}$$
 p->next->data = 5;



1 for 
$$(p = head; p - next != NULL; p = p - next);$$

$$p \rightarrow next \rightarrow data = 5;$$



3 p->next->data = 5;



;

$$p \rightarrow next \rightarrow data = 5;$$



,

$$p \rightarrow next \rightarrow data = 5;$$













```
1 for (p = head; p->next != NULL; p = p->next);
2 p->next = (listelem*)malloc(sizeof(listelem));
3 p->next->data = 5;
```

```
4 p->next->next = NULL;
```



```
1 for (p = head; p->next != NULL; p = p->next);
2 p->next = (listelem*)malloc(sizeof(listelem));
3 p->next->data = 5;
```

```
4 p->next->next = NULL;
```



If the list is empty, checking p->next != NULL is not possible, this case must be managed separately!

# Inserting element to the end of the list with a furet of the structure of the list with a furet of the structure of the struc

```
listelem *push_back(listelem *head, int d)
1
2
   ſ
3
     listelem *p;
4
     if (head == NULL) /* empty list should be
5
                managed separately */
6
       return push_front(head, d);
7
8
     for (p = head; p->next != NULL; p = p->next);
9
     p->next = (listelem*)malloc(sizeof(listelem));
10
   p->next->data = d;
11
     p->next->next = NULL;
12
     return head;
13
   }
14
                                                          link
   listelem *head = NULL;
1
2
   head = push_back(head, 2);
```



- If we have to traverse and process our data several times, it is worth sorting it
- Arrays:
  - re-locating a single element involves a lot of data movements
  - we fill up the array and order it afterwards



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  - we fill up the array and order it afterwards
- Lists:
  - re-locating a single element involves only the modification of pointers, the elements will remain at the same address in the memory
  - it is better to build up our list in a sorted way



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- The new element must be inserted before the first element that is larger then it
- In the present structure each element "can see" only behind itself, so we cannot insert element before another



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  - re-locating a single element involves a lot of data movements
  - we fill up the array and order it afterwards
- Lists:
  - re-locating a single element involves only the modification of pointers, the elements will remain at the same address in the memory
  - it is better to build up our list in a sorted way
- The new element must be inserted before the first element that is larger then it
- In the present structure each element "can see" only behind itself, so we cannot insert element before another
- We will use two pointers for traversing the list, one of them will be one step behind (delayed)
- We will insert after the delayed pointer

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```
1 q = head; p = q->next;
2 while (p != NULL && p->data <= data) { /* shortcut */
3 q = p; p = p->next;
4 }
5 q->next = (listelem*)malloc(sizeof(listelem));
6 q->next->data = 4;
7 q->next = p;
```


















































```
1 q = head; p = q->next;
2 while (p != NULL && p->data <= data) { /* shortcut */
3 q = p; p = p->next;
4 }
5 q->next = (listelem*)malloc(sizeof(listelem));
6 q->next->data = 4;
7 q->next->next = p;
```



# Inserting element into a sorted list with a functio

```
listelem *insert_sorted(listelem *head, int d)
1
   {
2
     listelem *p, *q;
3
4
      if (head == NULL || head->data > d) /* shortcut */
5
        return push_front(head, d);
6
7
      q = head;
8
     p = q - next;
9
     while (p != NULL && p->data <= d) /* shortcut */ {</pre>
10
        q = p; p = p - > next;
11
      }
12
13
     q->next = (listelem*)malloc(sizeof(listelem));
14
     q->next->data = d;
     q \rightarrow next \rightarrow next = p;
15
      return head;
16
   }
                                                                link
17
```









































This algorithm can be used only if we may modify the existing part of the list – others do not refer to it. But in many times this is not like that!

```
1 p = head;
2 while (p->next->next != NULL)
3 p = p->next;
4 free(p->next);
5 p->next = NULL;
```



```
1 p = head;
2 while (p->next->next != NULL)
3 p = p->next;
4 free(p->next);
5 p->next = NULL;
```





```
1 p = head;
2 while (p->next->next != NULL)
3 p = p->next;
4 free(p->next);
5 p->next = NULL;
```







```
1 p = head;
2 while (p->next->next != NULL)
3 p = p->next;
4 free(p->next);
5 p->next = NULL;
```







```
1 p = head;
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3 p = p->next;
4 free(p->next);
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4 free(p->next);
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2 while (p->next->next != NULL)
3 p = p->next;
4 free(p->next);
5 p->next = NULL;
```



```
1 p = head;
2 while (p->next->next != NULL)
3 p = p->next;
4 free(p->next);
5 p->next = NULL;
```



```
1 p = head;
2 while (p->next->next != NULL)
3 p = p->next;
4 free(p->next);
5 p->next = NULL;
```



If the list is empty or it contains only one element, the expression p->next->next doesn't make any sense.

# Deleting element from the end of the list with a trice in the second sec

```
listelem *pop_back(listelem *head)
1
   ł
2
     listelem *p;
3
4
     if (head == NULL ) /* empty */
5
6
       return head;
7
     if (head->next == NULL) /* only one element */
8
       return pop_front(head);
9
10
     for (p = head; p->next->next != NULL; p = p->next);
11
     free(p->next);
12
     p->next = NULL;
13
     return head;
14
                                                           link
15
   }
```


```
Deleting the data = 3 element
```

```
q = head; p = head->next;
1
  while (p != NULL && p->data != data) {
2
     q = p; p = p -> next;
3
  }
4
  if (p != NULL) { /* now we have it */
5
6
     q->next = p->next;
     free(p);
7
  }
8
```











```
q = head; p = head->next;
1
  while (p != NULL && p->data != data) {
2
     q = p; p = p -> next;
3
  }
4
  if (p != NULL) { /* now we have it */
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6
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     free(p);
7
  }
8
```





```
q = head; p = head->next;
1
  while (p != NULL && p->data != data) {
2
     q = p; p = p -> next;
3
  }
4
  if (p != NULL) { /* now we have it */
5
6
     q->next = p->next;
     free(p);
7
  }
8
```





```
q = head; p = head->next;
1
  while (p != NULL && p->data != data) {
2
     q = p; p = p -> next;
3
  }
4
  if (p != NULL) { /* now we have it */
5
6
     q->next = p->next;
     free(p);
7
  }
8
```







```
q = head; p = head->next;
1
  while (p != NULL && p->data != data) {
2
     q = p; p = p - > next;
3
  }
4
  if (p != NULL) { /* now we have it */
5
6
     q->next = p->next;
     free(p);
7
  }
8
```





```
Deleting the data = 3 element
```

```
q = head; p = head->next;
1
  while (p != NULL && p->data != data) {
2
     q = p; p = p - > next;
3
  }
4
  if (p != NULL) { /* now we have it */
5
6
     q->next = p->next;
     free(p);
7
8
  }
```







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#### Def Traversing Stack Insertion Deleting

# Deleting a given element from list





















```
q = head; p = head->next;
1
  while (p != NULL && p->data != data) {
2
     q = p; p = p -> next;
3
  }
4
  if (p != NULL) { /* now we have it */
5
6
     q->next = p->next;
     free(p);
7
8
  }
```











### Deleting an entire list



link

```
void dispose_list(listelem *head)
{
  while (head != NULL)
  head = pop_front(head);
}
```





 We have everything we need, but it was really cumbersome, because





- We have everything we need, but it was really cumbersome, because
  - we can insert element only after (behind) an element





- We have everything we need, but it was really cumbersome, because
  - we can insert element only after (behind) an element
  - we can delete only an element behind another element





- We have everything we need, but it was really cumbersome, because
  - we can insert element only after (behind) an element
  - we can delete only an element behind another element
  - empty lists and lists with only one element must be handled separately when inserting or deleting

### Thank you for your attention.