## Binary trees Basics of Programming 1



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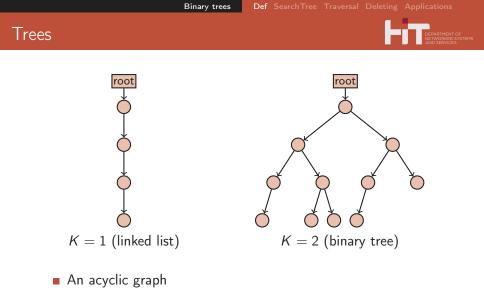
#### 1 Binary trees

- Definition
- Binary search trees

- Traversal
- Deleting
- Further applications

# Chapter 1

Binary trees

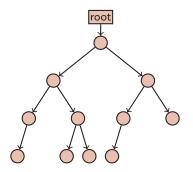


- Every node has exactly one incoming edge
- *K*-ary tree: every node has at most *K* outgoing edges

Binary trees

### Binary trees





Declaration of the binary tree data structure

```
1 typedef struct tree {
2   int data;
3   struct tree *left, *right;
4 } tree_elem, *tree_ptr;
```

Typically we typedef not only the struct, but also the pointer

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# 31 25 28 42 12 35 11 48 30 33 26

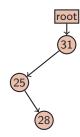
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- Sub-tree to the right: only elements greater than the node
- The structure of the tree depends on the insertion order of the elements!



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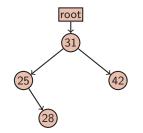




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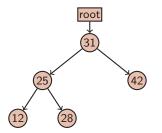




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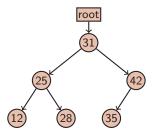




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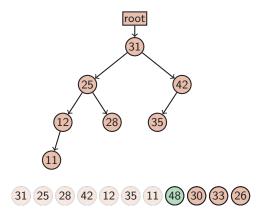




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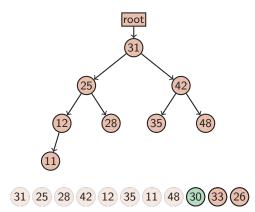
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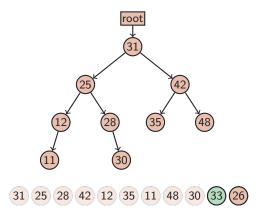
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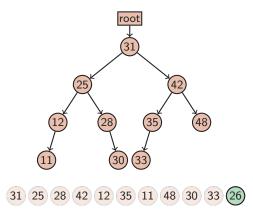
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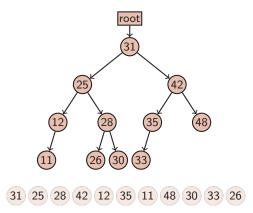
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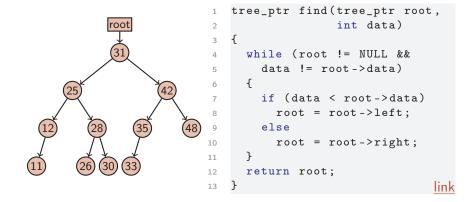
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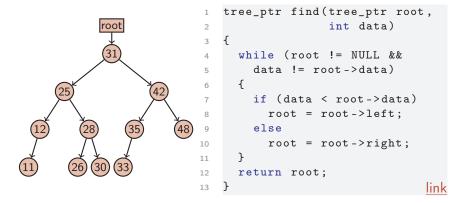


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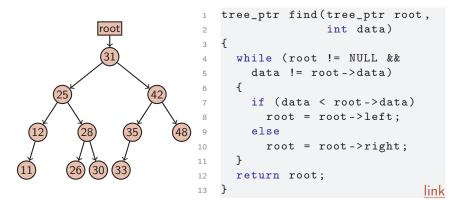






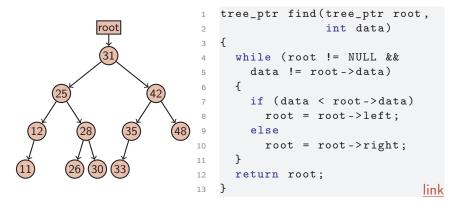
This is not recursive yet





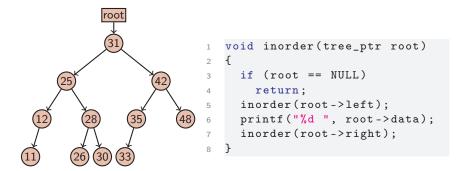
- This is not recursive yet
- In a depth-*d* tree the max. number of steps is *d*



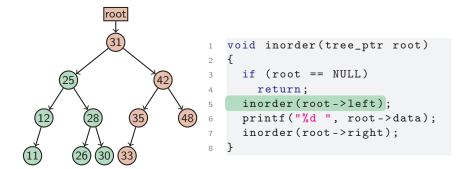


- This is not recursive yet
- In a depth-*d* tree the max. number of steps is *d*
- If the tree is balanced and has *n* elements  $\Rightarrow \approx \log_2 n$  steps!



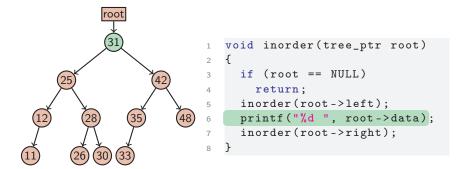






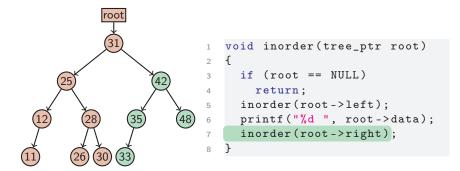
#### 11 12 25 26 28 30





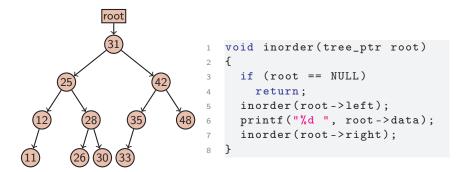
#### 11 12 25 26 28 30 31





#### 11 12 25 26 28 30 31 33 35 42 48





#### 11 12 25 26 28 30 31 33 35 42 48

in-order traversal



- 2 root element
- 3 right sub-tree

With this traversal the nodes are visited in increasing order of their values

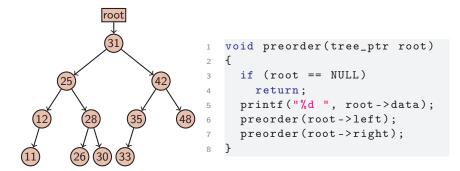


An other implementation of the in-order traversal:

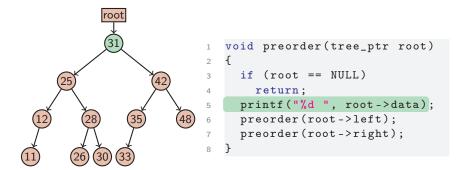
```
void inorder(tree_ptr root)
{
    if (root->left != NULL)
        inorder(root->left);
    printf("%d ", root->data);
    if (root->right != NULL)
        inorder(root->right);
    }
```

But in this case the caller has not make sure that root != NULL holds



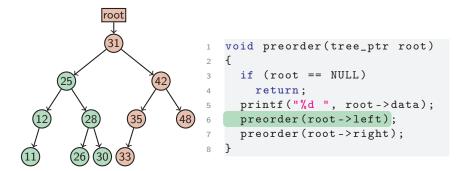






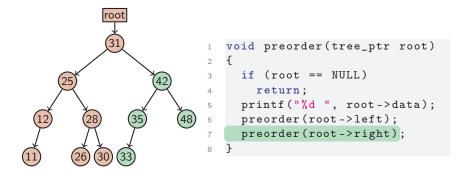
31





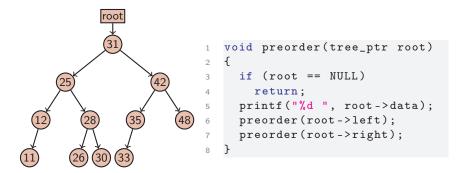
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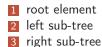
#### 31 25 12 11 28 26 30 42 35 33 48





#### 31 25 12 11 28 26 30 42 35 33 48

pre-order traversal



Saving the elements of the tree in this order, and building it again, the structure of the tree can be fully reconstructed.

### Building a tree



Inserting a new node to the tree

```
tree_ptr insert(tree_ptr root, int data)
1
2
   ł
     if (root == NULL) {
3
       root = (tree_ptr)malloc(sizeof(tree_elem));
4
       root->data = data;
5
     }
6
     else if (data < root->data)
7
       root->left = insert(root->left, data);
8
     else
9
       root->right = insert(root->right, data);
10
     return root;
11
   }
12
```

link

### Building a tree



Inserting a new node to the tree

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tree_ptr insert(tree_ptr root, int data)
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10
     return root;
11
12
   }
```

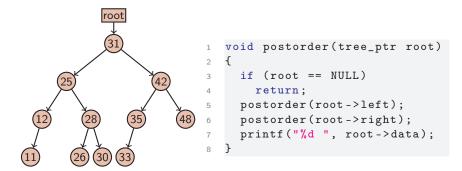
Usage of this function:

```
1 tree_ptr root = NULL;
2 root = insert(root, 2);
3 root = insert(root, 8);
4 ...
```

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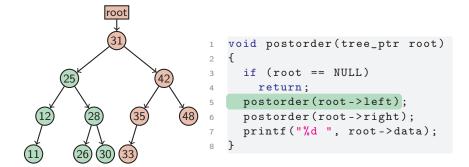
### Post-order traversal





#### Post-order traversal

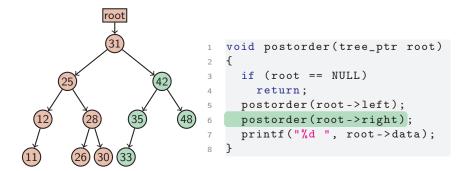




#### 11 12 26 30 28 25

#### Post-order traversal

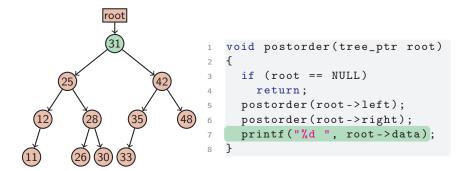




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#### Post-order traversal

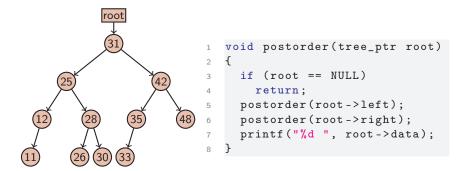




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#### Post-order traversal



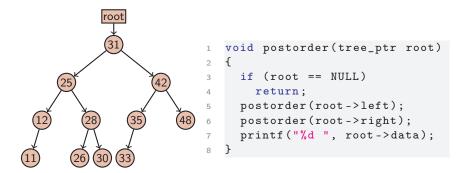


#### 11 12 26 30 28 25 33 35 48 42 31

- post-order traversal
  - left sub-tree
  - 2 right sub-tree
  - 3 root element

#### Post-order traversal





#### 11 12 26 30 28 25 33 35 48 42 31

- post-order traversal
  - left sub-tree
  - 2 right sub-tree
  - 3 root element

In this order the leaves of the tree are visited first  $\rightarrow$  application: releasing/deleting a tree

### Deleting a tree by post-order traversal

```
void delete(tree_ptr root)
  ł
2
    if (root == NULL) /* empty tree: nothing to delete *
3
       return;
4
    delete(root->left); /* post-order traversal */
5
    delete(root->right);
6
    free(root);
7
8
  }
                                                          link
```

### Deleting a tree by post-order traversal

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void delete(tree_ptr root)
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    if (root == NULL) /* empty tree: nothing to delete >
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4
    delete(root->left); /* post-order traversal */
5
    delete(root->right);
6
    free(root);
7
8
  }
                                                          link
```

A program segment (without memory leaks):

```
tree_ptr root = NULL;
root = insert(root, 2);
root = insert(root, 8);
...
delete(root);
root = NULL;
```



#### Write a recursive function (max. 10 lines), that



Write a recursive function (max. 10 lines), that

determines the depth of a tree



• Write a recursive function (max. 10 lines), that

- determines the depth of a tree
- calculates the count / the sum / the average of the values stored in the nodes of the tree



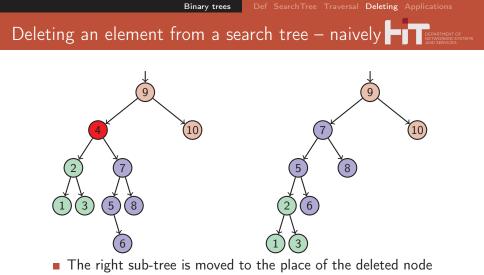
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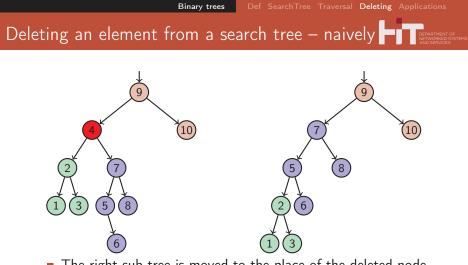
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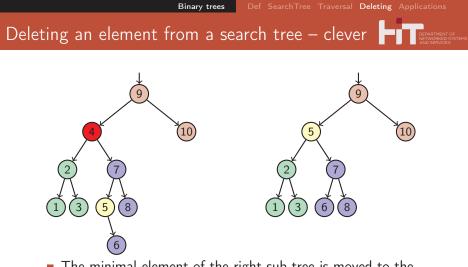
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  - calculates the count / the sum / the average of the values stored in the nodes of the tree
- Write a iterative function (max. 10 lines), that
  - computes the minimum and the maximum of the values stored in the nodes
  - returns the pointer to the node storing the maximal / minimal value of the tree



The left sub-tree is inserted to below the minimal element of the right sub-tree



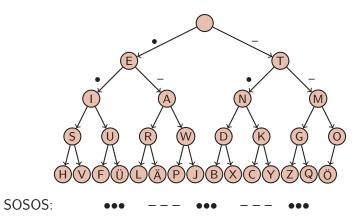
- The right sub-tree is moved to the place of the deleted node
- The left sub-tree is inserted to below the minimal element of the right sub-tree
- The tree is getting imbalanced!



- The minimal element of the right sub-tree is moved to the place of the deleted node
- This element could have only a right sub-tree, it is moved one level up, to its old place

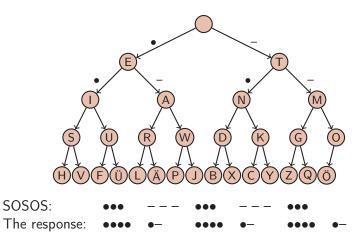
#### Morse decoding tree





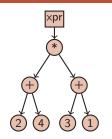
## Morse decoding tree





## Evaluating mathematical expressions

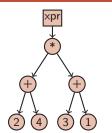




- Storing math expressions in a tree
- Leaves  $\rightarrow$  numeric constants
- $\blacksquare$  Branches  $\rightarrow$  two-operand operators
- In the example: (2+4) \* (3+1)

## Evaluating mathematical expressions





- Storing math expressions in a tree
- Leaves → numeric constants
- $\blacksquare$  Branches  $\rightarrow$  two-operand operators
- In the example: (2+4) \* (3+1)

```
int eval(tree_ptr xpr)
1
  ł
2
    char c = xpr->data;
3
    if (isdigit(c)) /* stopping condition */
4
      return c - '0';
5
    if (c == '+')
6
7
      return eval(xpr->left) + eval(xpr->right);
    if (c == '*')
8
      return eval(xpr->left) * eval(xpr->right);
9
```

link



Let us introduce variable x as a leaf node as well:

```
double feval(tree_ptr xpr, double x)
1
   ł
2
     char c = xpr->data;
3
     if (isdigit(c))
4
       return c - '0';
5
     if (c = 'x')
6
       return x;
7
     if (c == '+')
8
       return feval(xpr->left, x) + feval(xpr->right, x);
9
     if (c == '*')
10
       return feval(xpr->left, x) * feval(xpr->right, x);
11
   }
                                                            link
12
```

# Evaluating the derivative of a function



Let us take the derivative of the function! The rules are:

```
c' = 0
     x' = 1
     (f + g)' = f' + g'
     (f \cdot g)' = f' \cdot g + f \cdot g'
   double deval(tree_ptr xpr, double x)
1
   ł
2
3
     char c = xpr->data;
     if (isdigit(c)) /* stopping condition */
4
       return 0.0;
5
     if (c == 'x')
                         /* stopping condition */
6
       return 1.0:
7
     if (c == '+')
8
       return deval(xpr->left, x) + deval(xpr->right, x);
9
     if (c == '*')
10
       return deval(xpr->left, x) * feval(xpr->right, x) +
11
          feval(xpr->left, x) * deval(xpr->right, x);
12
                                                           link
13
   }
```

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Thank you for your attention.