









Computer Architectures

1. A card and code based door-lock

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- Designing the hardware
 - The building blocks of the system
 - Adding memory to the CPU
 - Adding peripherals to the CPU
- Designing the software
 - State-machine model of the system
 - Flow-charts of the algorithms
 - Implementation



THE HARDWARE

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A 8-bit CPU

- Data bus: 8 bit wide
- Address bus: 16 bit wide
- Multiplexed memory and I/O buses
- RISC instruction set
- 3-operand instruction format
- Starting address: 0000h
- Interrupt subroutine address: 1000h

- Memory
 - ROM: to store instructions and program constants
 - RAM: to store program variables and stack
- Let us use 8kB ROM and 8kB RAM
 - 8 bit data bus
 - 13 bit address bus



- Card reader
 - OE: output enable
 - It has an 8 bit output: D
 - If $\overline{S}/D = 0$: it gives back an 1, if a new card has arrived
 - If $\overline{S}/D = 1$: it gives back the code (ID) of the last card
 - INT: generates an interrupt, if a new card is recognized



- Keypad
 - OE: output enable
 - It has a 8 bit output D:
 - If $\overline{S}/D = 0$: it gives back an 1 if a button has been pressed
 - If $\overline{S}/D = 1$: it gives back the code of the last key
 - INT: generates an interrupt upon each key press



- Door lock
 - CS: chip select
 - WR/RD
 - If $\overline{WR}/RD = 1$: sets L according to the state of the door
 - (open=1, locked=0)
 - If $\overline{WR}/RD = 0$: opens the door. The door automatically locks again after a timeout.



- LED lights
 - Indicate the status of the door (red/green)
- Additional elements:
 - Decoder for interfacing the memory
 - Comparators for detecting the addresses of peripherals
 - D flip-flop for driving the LED lights

Adding memory to the CPU

- We have 8 kB ROM + 8 kB RAM
- Memory map:

BME

E000h-FFFFh					
C000h-DFFFh					
A000h-BFFFh					
8000h-9FFFh					
6000h-7FFFh		A15			A ₀
4000h-5FFFh		010 1 010 0	1111 0000	1111 0000	1111 0000
		001 1	1111	1111	1111
2000h-3FFFh	RAM	001 0	0000	0000	0000
		0001	1111	1111	1111
0000h-1FFFh	ROM	0000	0000	0000	0000

- A13-A15 determine which module to use
- A₀-A₁₂ determine the byte position in the selected module

Adding memory to the CPU



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Adding the card reader



Adding the keypad: the same, with base address 3Eh

Adding the door lock



Adding the LED lights



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THE SOFTWARE

State-machine of the system

ST (state) = number of correct code digits typed so far on the keypad



Transitions of the state-machine



- Card event: ST=0, and store the card ID in a local variable
- Correct digit: (if ST!=255 then) ST=ST+1,
 - if now ST==4 then: ST=255 and open door
- Bad digit: ST=255

BME



- The array of 4-digit key codes for each card ID
 - The card ID is 8 bit wide \rightarrow 256 different IDs are possible
 - An array is used
 - with 256 entries, one for each card ID
 - each entry is the correct code (4-byte long) of the given card
 - Total memory consumption: 256*4=1024 bytes (=400h)
 - The *i*th code digit of card *j* is located at
 - array start address + *j**4 + *i* (*j* in {0..255}, *i* in {0..3})
- Local variables:
 - 1-byte integer variable ST
 - 1-byte integer variable CARDID



- The program has two parts:
 - Main program
 - Interrupt service routine
- The purpose of the main program:
 - Monitors the state of the door
 - Adjusts the LEDs accordingly
- The purpose of the interrupt service routine
 - Handles card events and key presses
 - Opens the door, if the correct code is given



- Initialization:
 - Set stack pointer
 - Set the value of ST (state variable)
 - Enable interrupts
- In an infinite loop:
 - Ask the door for the state
 - Set the LED-s accordingly



The interrupt service routine

- First, we have to find out the interrupt source (polling!)
- Then, we apply the transition on the state machine







The interrupt service routine



The interrupt service routine

The end of interrupt service routine:



Placement decisions

Content of the RAM

- 1-byte integer variable ST: **2000h** (first byte of the RAM)
- 1-byte integer variable CARDID: 2001h (second byte of the RAM)
- Initial stack pointer: **3FFFh**, the highest address of the RAM, since it grows downwards

Content of the ROM

- The program, starting at **0000h**
- The interrupt service routine, starting at **1000h**
- The array of correct codes for each card
 - Size: 400h
 - Can be placed anywhere, where space is available
 - Let us put it to 500h
 - The main program is small, it does not go beyond 500h
 - The end of the array is **8FFh**, which is below the interrupt routine

Implementation of the main program:

	ORG 0000h	start code at this address
	SP – 3FFFh	set stack pointer
	MEM[2000h] ~ 255	ST=255 (Initial state)
	EI	enable interrupts
label:	R0 ← IO[4Ch]	read door status
	IO[5Dh] ← RØ	update LEDs
	JUMP label	jump back to reading

```
    Array of 4-byte codes for each card
ORG 500h
codes: DB 1, 3, 4, 7
DB 5, 7, 2, 9
DB 8, 2, 0, 8
DB 3, 1, 8, 9
```

BME

Implementation of the interrupt service routine:

ORG 1000h CPU jumps at this address on interrupt PUSH RØ save R0/R1/R2 register to stack PUSH R1 PUSH R2 **R0** ← **IO**[2Eh] read card reader state JUMP keycode IF R0==0 jump, if IT source is not the reader $R0 \leftarrow I0[2Fh]$ read card ID MEM[2001h] ~ R0 save to variable CARDID MEM[2000h] ~ 0 ST=0 JUMP end jump to the end of interrupt routine

Implementation of the interrupt service routine:

BME

keycode: R0 ← IO[3Fh] R0=new digit $R1 \leftarrow MEM[2000h]$ R1=ST JUMP end IF R1==255 we are not expecting a digit **R2** ← MEM[2001h] R2=CARDID $R2 \leftarrow R2 * 4$ R2=4*CARDID $R2 \leftarrow R1 + R2$ R2=4*CARDID+ST R2 ← MEM[R2+codes] R2=the correct digit JUMP match IF R2==R0 Jump if new digit is correct MEM[2000h] ~ 255 state machine: ST=225 (init. state) JUMP end and go to the end of interrupt match: **R1** ← **R1+1** If correct, increment ST MEM[2000h] ~ R1 and save ST to memory JUMP end IF R1<4 If this is not the last digit, go to the end **IO**[4Ch] ~ 1 If last digit, open the door MEM[2000h] ~ 255 state machine: ST=225 (init. state) JUMP end Go to the end of interrupt

Implementation of the interrupt service routine:

end:	POP	R2	restore registers in a reverse order
	POP	R1	
	POP	R0	
	EI		enable further interrupts
	RET		return from interrupt



- Can we solve the problem without RAM?
 - Yes.
 - But we lose the stack
 - No more interrupts!!!
 - Main program continuously polls:
 - The state of the door (to update LEDs)
 - The state of the card reader
 - The state of the keypad
 - We can not put ST and CARDID into the RAM
 - We can store them in registers permanently