Checking the Accuracy of Siitperf

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Abstract: Siitperf is the world’s first free software RFC 8219 compliant SIIT (also called stateless NAT64) tester, which implements throughput, frame loss rate, latency and packet delay variation tests. In this paper, we show that the reliability of its results mainly depends on the accuracy of the timing of its frame sender algorithm. We also investigate the effect of Ethernet flow control on the measurement results. Siitperf is calibrated by the comparison of its results with that of a commercial network performance tester, when both of them are used for determining the throughput of the IPv4 routing of the Linux kernel.

Keywords: siitperf, accuracy, calibration

1. Introduction

RFC 8219 [1] has defined a benchmarking methodology for the high number of IPv6 transition technologies [2] by classifying them into a small number of categories and defining benchmarking procedures for each category. As far as we know, our siitperf [3] is the world’s first free software RFC 8219 compliant SIIT [4] (also called stateless NAT64) tester, written in C++ using DPDK (Data Plane Development Kit) [5] available from GitHub [6]. Being a measurement tool, the accuracy of siitperf is a key issue, which we examine in this paper. To that end, first, we give a very short introduction to siitperf only up to the measure necessary to understand the rest of this paper. Then, we define our error model by overviewing the most important factors that could cause unreliable measurement results. Next, we examine the effect of Ethernet flow control on the measurement results. After that, we measure the throughput of the same DUT (Device Under Test) using a commercial network performance tester and siitperf and compare their results. Finally, we discuss our results and disclose our plans for further research.
2. A short introduction to \texttt{siitperf}

In order to provide the reader with the necessary background information for the understanding of the rest of this paper, we give a short overview of \texttt{siitperf} on the basis of our open access paper [3], in which all the details can be found. From among the measurement procedures of RFC 8219, \texttt{siitperf} currently supports the following ones: \textit{throughput}, \textit{frame loss rate}, \textit{latency} and \textit{packet delay variation}. All these procedures require the sending and receiving of certain IPv4 or IPv6 frames (more precisely: Ethernet frames containing IPv4 or IPv6 packets) at a pre-defined constant frame rate according to the test setup shown in Fig. 1, where the \textit{Tester} device benchmarks the \textit{DUT} (Device Under Test). Although the arrows would imply unidirectional traffic, testing with bidirectional traffic is required by RFC 8219 and testing with unidirectional traffic is optional. Of course, both X and Y in IPvX and IPvY are from the set of \{4, 6\}. Naturally, if we are talking about SIIT (stateless NAT64), then it implies that X≠Y. However, \texttt{siitperf} allows X=Y, that is, it can also be used for benchmarking an IPv4 or IPv6 router.

The first two of the supported benchmarking procedures (\textit{throughput} and \textit{frame loss rate}) require only the above mentioned sending of test frames at a constant rate and counting of the received test frames. The core measurement of both procedures is the same. The difference is that throughput measurement requires to find the highest rate at which the DUT can forward all the frames without loss, whereas the frame loss rate measurement requires to perform the core measurement at various frame rates to determine the frame loss rate at those specific frame rates. The core measurement of both tests is implemented in \texttt{siitperf-tp} and two different benchmarking procedures are performed by two different shell scripts. It is also very important that RFC 8219 requires the execution of the throughput measurement (that is the binary search to determine the highest lossless frame rate called as \textit{throughput}) at least 20 times, then the median of the results is calculated as the final result, and the first percentile and the 99th percentile of the results are used as indices of dispersion.

\begin{center}
\includegraphics[width=0.5\textwidth]{figure1.png}
\end{center}

\textit{Figure 1. Single DUT test setup [1]}


The latency benchmarking procedure is performed at the rate determined by the throughput measurement. RFC 8219 requires the tagging of at least 500 frames, which are used for latency measurement. Latency measurement requires that timestamps are stored immediately after sending and receiving of tagged frames. The latency for each tagged frame is calculated as the difference of the receiving and sending timestamps of the given frame. (For the details of the calculation of the final result, please refer to RFC 8219.) The latency benchmarking procedure is implemented by `siitperf-lat`, which is an extension of `siitperf-tp`.

From our point of view, the packet delay variation benchmarking procedure is similar to the latency benchmarking procedure, but it requires timestamping of every single frame. (The final result calculation is different, but now we do not deal with that.) The packet delay variation benchmarking procedure is implemented by `siitperf-pdv`, which is also an extension of `siitperf-tp`.

Internally, `siitperf` uses `TSC` (Time Stamp Counter) for time measurements, which is a very accurate and computationally inexpensive solution (it is a CPU register, which can be read by a single CPU instruction: `RDTSC` [7]).

To achieve as high performance as possible, all test frames used by `siitperf-tp` and `siitperf-lat` are pre-generated (including the tagged frames). The test frames of `siitperf-pdv` are prepared right before sending by modifying a set of pre-generated frames: their individual identifiers and checksums are rewritten.

Regarding our error model, it is important that the sending and receiving of the frames are implemented by sender and receiver functions, which are executed as threads by the CPU cores specified by the user in the configuration file.

3. Our error model

3.1. Accuracy of the timing of frame sending

There is an excellent paper that examines the accuracy of the timing of different software packet generators [8]. It points out that the inter-sending time of the packets is rather imprecise at demanding frame rates, if pure software methods are used. It also mentions the buffering of the frames by the NIC (Network Interface Card) among the root causes of this phenomenon, what we have also experienced and reported: our experience was that when a packet was reported by the DPDK function as “sent”, it was still in the buffer of the NIC [3]. Unfortunately, this buffering completely discredits any investigation based on using timestamps stored at the sending of the frames. (Even if we store timestamps both before and after the sending of a frame, we may not be sure, when the frame was actually sent.)
Imprecise timing may come from various root causes. At demanding frame rates, one of them is that our contemporary CPUs use several solutions to increase their performance including caching, branch prediction, etc. and they usually provide their optimum performance only after the first execution (or after the first few executions) of the core of the packet sending cycle, thus the first (few) longer than required inter-sending time(s) is/are followed by shorter ones to compensate the latency. This compensation depends on the timing algorithm of the sender function. For example, the original implementation of dns64perf++ used a sophisticated algorithm that intended to distribute the compensation of such initial latency for the rest of the measurement time [9]. Unfortunately, the compensation algorithm did not work well and thus the sending rate was somewhat lower than required from the beginning of the measurement for a long time, and it was significantly higher than required at the end [10]. Therefore, we have replaced the timing algorithm with a simpler one that promptly compensates the latency [10]. We followed the same approach in siitperf, thus it sends the test frame (if it can), when its time has arrived, with no respect to what has happened before [3]. Therefore, siitperf very likely produces micro burst(s) at rates close to the upper limit of its sending performance.

Unfortunately, we did not have a NetFPGA device used by the authors of [8], therefore, we decided to check, how the imprecise timing of siitperf influences its measurement results. Our error model is that traffic with not exact inter-arrival time may have the following influence on the throughput test results:

1. The median decreases, because the imprecise timing causes sometimes overload and thus frame loss at lower rates than the throughput rate with precise timing.

2. The dispersion of the results increases, because some random events (like interrupts) influence each execution of the test differently.

The actual frame loss caused by the imprecise timing may also depend on a further parameter, namely, if Ethernet flow control (IEEE 802.3x) is used or not, because flow control may “iron out” the random peaks of the frame sending rate caused by imprecise timing.

Therefore, first, we test how the presence or absence of flow control influences the results. This phenomenon is interesting by itself, and the results of this comparison proved to be very important due to the limitations of our next examination.

Then, we benchmark the same DUT with both a calibrated tester and siitperf so that we can see the difference. The fact that siitperf is able to perform pure
IPv4 or IPv6 benchmarking tests, allowed us to use an RFC 2544 [11] compliant legacy tester. This solution has also its limitations: although RFC 8219 has taken the throughput and frame loss rate tests verbatim from RFC 2544, the latency test has been redefined (it requires at least 500 tagged frames instead of a single one) and packet delay variation measurement is a completely new one. Thus they cannot be validated by an RFC 2544 tester.

We note that even if we can directly check the accuracy of frame sending of \texttt{siitperf-tp} only, we expect that the accuracy of frame sending of the other two programs is not worse, either. As for \texttt{siitperf-lat}, the relatively low number of tagged frames, which are distributed evenly, cannot make any significant effect. As for \texttt{siitperf-pdv}, the setting of their individual identifier and checksum requires some time, and thus there is non-zero lower bound for their inter-frame time, at least in theory. We note that it guarantees nothing in practice due to the fact of NIC buffering: back-to-back frames (that is frames with minimum inter-frame gap) may still occur.

\textbf{3.2. Consideration of other errors}

Unlike the sender function that sends frames individually, the receiver function may receive multiple frames together to ensure high performance. This can surely not cause any problem with the throughput and frame loss rate measurements, because the frames are only counted. The receiving timestamps of latency and packed delay variation tests may be influenced, but they are also influenced by buffering even if they are taken out from the receive buffer individually.

The sending and receiving timestamps are subject to further errors due to the fact that interrupts may occur between the sending/receiving of the frames and taking the timestamp by the execution of the RDTSC machine code instruction. This is a kind of error we cannot measure. As for latency measurements, one possible mitigation can be, if the user sets the number of time stamps to be used to a significantly higher value than the required minimum 500 (\texttt{siitperf} supports up to 50,000) and thus the calculation of the 99.9th percentile removes the errors, if they are rare enough. This mitigation automatically applies for packet delay variation tests, as all frames are time stamped.

Although it is the responsibility of the user to specify the four cores that execute the sending and receiving threads so that they belong to the same physical CPU as the main core (used for starting the program), \texttt{siitperf} does some sanity checks if the TSC-s of the four CPU cores are synchronized with that of the main core. (Otherwise the TSC values specified for starting and stopping the experiment as well
as the differences of the timestamps of the corresponding senders and receivers would be invalid.)

We believe that all other errors including the conversions between (milli)seconds and TSC, the counting of the sent and received frames, the calculations with the timestamps, etc. are subject to general software testing and verification procedures.

4. Investigation of the effect of the Ethernet flow control

To be able to investigate, how the presence or the absence of the Ethernet flow control influences the results, we needed a test system that is free from any other effects that may make our results noisy. Based on our SIIT benchmarking experience [12], we have chosen to reuse a previously built tests system, which was build up by two identical Dell PowerEdge C6220 servers in the NICT StarBED, Japan. The very same system was also used for benchmarking the extension of siitperf with the ability of using random source and destination port numbers [13] as required by RFC 4814 [14]. We have taken the following description of the test system from that paper [13]. “The servers were equipped with two 2GHz Intel Xeon E5-2650 CPUs having 8 cores each, 128GB 1333MHz DDR3 RAM and Intel 10G dual port X520 Ethernet network adapters. The Debian Linux operating system was updated to version 9.13 on all computers. The Linux kernel version was: 4.9.0-4-amd64. The DPDK version was 16.11.11-1+deb9u2.” [13]

![Diagram](image)

*Figure 2. Measurement setup for IPv4 Linux kernel routing: throughput tests with and without Ethernet flow control.*
The “varport” branch of \textit{siitperf} was used, its latest commit was bfdddb5f on Aug 23, 2020.

We expected that the difference between the results with and without Ethernet flow control depends on the frame rate and we also wanted to test this hypothesis.

To achieve high enough frame rates, first, we bench marked IPv4 kernel routing using random source and destination port numbers as we did in [13]. The topology of the test system is shown in Fig. 2. The CPU clock rate was set to fixed 2GHz on both computers (using the same BIOS settings as specified in the appendix of [15]). All cores of the second CPU of the DUT were switched off using the \texttt{maxcpus=8} kernel parameter to avoid NUMA issues (please refer to [12] for a detailed explanation).

As for frames sizes, we used 64, 128 and 256 bytes from among the standard frame sizes recommended by RFC 8219, because the throughput of the DUT was limited by the performance of the DUT with these frame sizes, whereas throughput was limited by the maximum frame rate of the 10G Ethernet for all higher standard frame sizes.

Technical note: \textit{siitperf} interprets the specified frame size values for IPv6 frames and uses 20 bytes less for IPv4 frames, therefore we always set 20 bytes higher values. (This is important, if someone would like to repeat our experiments.)

As required by RFC 8219, we used bidirectional traffic and 60s long tests at each step of the binary search. The “Error” parameter of the binary search was set to 1. (Error means that the binary search may stop, when: $upper \_ limit - lower \_ limit \leq Error$.) The results are shown it Table 1. Besides the values required by RFC 8219 (median, first percentile and 99th percentile of the 20 throughput results for each frame size), we have also calculated another value to express the consistent or scattered nature of the results, which we named \textit{dispersion} and defined as follows:

Table 1. IPv4 Linux kernel routing performance with and without flow control, Dell PowerEdge C6220 servers, fixed 2GHz CPU clock rate, 8 active CPU cores, RFC 4814 random port numbers

<table>
<thead>
<tr>
<th>frame size</th>
<th>with flow control</th>
<th>without flow control</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>64</td>
<td>128</td>
</tr>
<tr>
<td>median</td>
<td>3,432,658</td>
<td>3,352,378</td>
</tr>
<tr>
<td></td>
<td>3,411,322</td>
<td>3,344,630</td>
</tr>
<tr>
<td>1st perc.</td>
<td>3,420,774</td>
<td>3,347,624</td>
</tr>
<tr>
<td></td>
<td>3,374,999</td>
<td>3,312,499</td>
</tr>
<tr>
<td>99th perc.</td>
<td>3,441,407</td>
<td>3,359,921</td>
</tr>
<tr>
<td></td>
<td>3,418,731</td>
<td>3,351,578</td>
</tr>
<tr>
<td>disp. (%)</td>
<td>0.60</td>
<td>0.37</td>
</tr>
<tr>
<td></td>
<td>1.28</td>
<td>1.17</td>
</tr>
</tbody>
</table>
\[
\text{Dispersion} = \frac{99^{th} \text{ percentile} - 1^{st} \text{ percentile}}{\text{median}} \cdot 100 \%
\]

We note that commercial Testers like the one we used in the next section, usually report the number of \textit{all} frames per second (including frames in both directions), but \texttt{siitperf} reports the number of frames per second \textit{per direction}. Thus the number of \textit{all} frames per second forwarded by the Linux kernel was the double of the numbers shown in Table 1.

Let us compare the results with and without flow control for each frame size individually. As for 64-byte frames, the median throughput with flow control (3,432,658fps) is about 0.6% higher than the median throughput without flow control (3,411,322fps). The lack of flow control has also significantly increased the dispersion of the results (from 0.6% to 1.28%). As for 128-byte frames, the median throughput with flow control (3,352,378fps) is only about 0.2% higher than the median throughput without flow control (3,444,630fps). Finally, with 256-byte frames, the median throughput with flow control (3,153,894fps) is only about 0.03% higher than the median throughput without flow control (3,152,872fps). We can observe that the difference between the results with and without flow control definitely decreases with the increase of the frame size.

We are satisfied with the results in the sense that the 3.4Mfps is more than the half of the 6.3Mfps maximum frame rate \texttt{siitperf} can achieve on the given hardware [13] and our results with and without Ethernet flow control are quite close to each other (even the largest difference is below 1%).

To be able to test the effect of the Ethernet flow control at a significantly lower frame rate using the very same test system, we used fixed port numbers. In this case, the packet processing at the DUT was not hashed to all 8 active CPU cores of the DUT, but only two CPU cores were used (one core per direction). (We could observe only the load caused by software interrupts.) This time we performed the throughput measurements using all standard frame sizes, as throughput was always lower than the theoretical maximum value for the media at the given frame size. Our results with flow control are shown in Table 2. As we expected, the dispersion of the results is lower than 1% for all frame rates. The median throughput is slightly decreases, when we increase to frame size from 64 bytes (885,643fps) through 128 bytes (878,256fps) to 256 bytes (857,575fps). There is significant decrease at 512 bytes (779,410fps) and median throughput remains constant (within measurement error) at all higher frame sizes. The investigation of the decrease of the median at 512 bytes is beyond the scope of our current paper, we just mention that it can also be observed in Fig. 3 of [12].
Our results without flow control are shown in Table 3. The difference of the median throughput between the results with flow control (885,643fps) and without flow control (880,381fps) is about 0.6% at 64 bytes frame size. Although this difference decreases to 0.36% at 128 bytes, but it is about 0.8% at 256 bytes frame size. Thus the increase of the frames size was not enough to make the difference diminish. For the following three standard frame sizes, this difference is about: 0.14%, 0.08%, 0.2%, and for the last two frame sizes, the difference is deliberately less than measurement error. Unfortunately, the dispersion of the results of the measurements without flow control is rather high: it exceeds 15% at 128 bytes frame size. At this point, we cannot tell whether this high dispersion is caused by the improper timing of sitiperf or by the nature of the DUT.

5. Calibration with a standard tester

We have built two tests systems to determine the IPv4 routing performance of the same DUT, which was a Sun Fire X4150 server with two Quad Core 2.83GHz Intel Xeon E5440 CPUs, four 2GB 667MHz DDR2 SDRAM modules and four Gigabit Ethernet ports. Debian 9.11 GNU/Linux operating system with 4.9.0-5-amd64
kernel was installed on it. The clock frequency of all 8 CPU cores was set to fixed 2.833GHz using the `cpufreq-set` command of the `cpufrequtils` package.

5.1. Reference measurements

To provide reference, the throughput of IPv4 Linux kernel routing was measured using a commercial Anritsu MP1590B Network Performance Tester. It had a four port Anritsu MU210212A 10/100/1000M Ethernet Module, and we used Port1 and Port2 of the module. The measurement setup is shown in Fig. 3.

As RFC 8219 has somewhat extended the standard frame sizes to be used for benchmarking originally defined in RFC 2544, we have chosen custom frame sizes and defined the following frame sizes: 64, 128, 256, 512, 768, 1024, 1280, 1518.

As required by RFC 8219, bidirectional traffic was used and full 60s length trials were executed and the “Loss Tolerance” parameter was set to 0%.

The Anritsu tester has a parameter called “Resolution”, which can be specified as the percentage of maximum frame rate of the media. Its smallest possible value is 0.01. As the theoretical maximum frame rate for Gigabit Ethernet with 64 byte frame size is 1,488,095, this setting means that the resolution (or with other words error) of the binary search was about 149fps (calculated as: 1,488,095/10,000).

We note that the Anritsu tester does not perform a binary search, if the test at the “Maximum Frame Rate” is successful. (We set the value of this parameter to 100%.)

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![Figure 3. Measurement setup for IPv4 Linux kernel routing: reference throughput test with a commercial Tester](image-url)
By default, flow control was not enabled on the Anritsu tester. When we enabled flow control, the Anritsu tester became practically unusable for throughput measurements, because it qualified all tests as successful, even if they lasted much longer than 60s. Thus, we could use this tester for meaningful measurements only, when flow control was disabled.

5.2. Measurements with siitperf

The parameters of the DUT were the same as in the previous case, but this time the Tester was a Dell PowerEdge R620 server with two six core 2GHz Intel Xeon E5-2620 CPUs, two 16GB 1600MHz DDR3 SDRAM modules and with an additional Intel I350-T4 Ethernet Server Adapter (needed for DPDK). Debian 9.11 GNU/Linux operating system with 4.9.0-11-amd64 kernel was installed on it. The version of DPDK was 16.11.9-1-deb9u2. As siitperf does not have version numbers yet, we can identify its version with its latest commit number 05247a1 on Jul 1, 2020. (This time, we used the master branch.)

The measurement setup is shown in Fig. 4. We used the same standard frame sizes mentioned before.

As required by RFC 8219, bidirectional traffic was used and full 60s length trials were executed, and the “Error” of the binary search was set to 1.

We note that as the binary-rate-alg.sh script distributed with siitperf supports only tests for a single pre-set frame size, with a single pre-set upper bound,
we have added a for cycle to the script with the appropriate frame sizes and the following upper bounds for the consecutive standard frame sizes: 1,500,000 850,000 460,000 240,000 160,000 120,000 100,000 82,000. They are wilfully somewhat higher than the theoretical maximum frame rates for the media with the given frame size, because we wanted to test and demonstrate how sitperf behaves, when the maximum frame rate for the media is achieved by the DUT. Our script performed the binary search for all standard frame sizes starting in the interval of 0 (as lower bound) and the above mentioned upper bound values.

Unlike with the Anritsu Tester, the log file of the DUT showed that flow control was enabled on the interfaces used for testing (Flow Control: Rx/Tx). We tried to switch off flow control using the same command as with the 10G Ethernet interfaces (ethtool -A interface rx off tx off), which has been executed without any error message, however flow control remained enabled.

5.3. Results

The Anritsu Tester reported the results in a form of a graph, which we include in Fig. 5 to facilitate an easy overview of the results. Except for the first two frame sizes, the throughput achieved its theoretical maximum value. However, this reporting format covers some very important details by displaying only the average

![Figure 5. IPv4 Linux kernel routing performance of the Sun server: reported by the Anritsu Tester with no flow control](image-url)
value of the measurement results. Therefore, we processed the detailed result file and calculated the median, first percentile, and 99th percentile of the 20 throughput results for each frame size, as required by RFC 8219. Please see our results in Table 4. We note that the Anritsu Tester reported the number of all frames per second (including frames in both directions), but we divided the results by two to report the number of frames per second per direction. We did so to show values comparable with the theoretical maximum frame rates given in Appendix A.1 of RFC 5180 [16].

One of the most conspicuous things in the table is the high dispersion of the results at 64-byte frame size. It is caused by a single outlier. We have investigated the case in the measurement log file, and we found that 8 frames were missing during that step of the binary search when the target rate was 34.37%. Of course, it meant that the test failed. After that, all tests were successful and thus the final result was 34.36%. This single outlier does not influence the median, but it is reflected by the first percentile and, therefore, in the dispersion, too.

As for the results at 128-byte frame size, the first percentile and the 99th percentile are nearly symmetrical around the median.

As for the results at 256-byte frame size, a single test failed at 100% due to the loss of a few frames, therefore, binary search was performed, which finished at 99.99%. All other tests passed at 100% and thus no binary search was performed.

Table 4. IPv4 Linux kernel routing performance of the Sun server: measured by the Anritsu Tester without flow control

<table>
<thead>
<tr>
<th>Fr. size</th>
<th>64</th>
<th>128</th>
<th>256</th>
<th>512</th>
<th>768</th>
<th>1024</th>
<th>1280</th>
<th>1518</th>
</tr>
</thead>
<tbody>
<tr>
<td>median</td>
<td>548,958</td>
<td>526,351</td>
<td>452,899</td>
<td>234,962</td>
<td>158,629</td>
<td>119,732</td>
<td>96,154</td>
<td>81,274</td>
</tr>
<tr>
<td>1st perc.</td>
<td>511,310</td>
<td>522,720</td>
<td>452,853</td>
<td>234,962</td>
<td>158,629</td>
<td>119,732</td>
<td>96,154</td>
<td>81,274</td>
</tr>
<tr>
<td>99th perc</td>
<td>553,720</td>
<td>529,561</td>
<td>452,899</td>
<td>234,962</td>
<td>158,629</td>
<td>119,732</td>
<td>96,154</td>
<td>81,274</td>
</tr>
<tr>
<td>disp. (%)</td>
<td>7.73</td>
<td>1.30</td>
<td>0.01</td>
<td>0.00</td>
<td>0.00</td>
<td>0.00</td>
<td>0.00</td>
<td>0.00</td>
</tr>
</tbody>
</table>

Table 5. IPv4 Linux kernel routing performance of the Sun server: measured by \textit{siitperf} with flow control

<table>
<thead>
<tr>
<th>Fr. size</th>
<th>64</th>
<th>128</th>
<th>256</th>
<th>512</th>
<th>768</th>
<th>1024</th>
<th>1280</th>
<th>1518</th>
</tr>
</thead>
<tbody>
<tr>
<td>median</td>
<td>549,297</td>
<td>522,413</td>
<td>452,930</td>
<td>234,986</td>
<td>158,652</td>
<td>119,752</td>
<td>96,173</td>
<td>81,294</td>
</tr>
<tr>
<td>1st perc.</td>
<td>547,850</td>
<td>521,285</td>
<td>452,853</td>
<td>234,986</td>
<td>158,650</td>
<td>119,752</td>
<td>96,173</td>
<td>81,294</td>
</tr>
<tr>
<td>99th perc</td>
<td>550,782</td>
<td>524,610</td>
<td>452,960</td>
<td>234,990</td>
<td>158,667</td>
<td>119,767</td>
<td>96,178</td>
<td>81,301</td>
</tr>
<tr>
<td>disp. (%)</td>
<td>0.53</td>
<td>0.64</td>
<td>0.01</td>
<td>0.00</td>
<td>0.01</td>
<td>0.01</td>
<td>0.01</td>
<td>0.01</td>
</tr>
</tbody>
</table>
No binary search was performed at any higher frame sizes, this is why their first percentiles and the 99th percentiles are equal with their medians.

The results of the throughput measurements with \texttt{siitperf} are shown in Table 5. The dispersion of the results is always below 1%, and it is practically 0 upwards from 256 bytes frame size, as the maximum frame rate for the media has limited the throughput. As the upper limit was set higher than the maximum frame rate for the media, \texttt{siitperf} executed binary search and it measured slightly higher values. It was possible for at least two reasons:

- As Appendix A.1 of RFC 5180 states: “Ethernet's maximum frame rates are subject to variances due to clock slop. The listed rates are theoretical maximums, and actual tests should account for a +/- 100 ppm tolerance.”
- The “TOLERANCE” parameter of \texttt{siitperf} was set to 1.00001, which means that 0.001% more time is allowed for sending.

There are two throughput values that were limited by the CPU performance: throughput measured with 64 bytes and 128 bytes frame sizes. The differences are 0.06% and 0.75%, which we consider good and acceptable, respectively.

6. Discussion and plans for future research

Our conditions for calibrating \texttt{siitperf} with a standard tester were far from ideal. We cannot tell the maximum frame rate, at which the CPU of the Dell PowerEdge R620 server would be able to generate frames, but it is very likely several million frames per second, thus the measured throughput around 550,000 fps was not at all close to it. The technical issue that we could use the Anritsu tester only without flow control, whereas we could use \texttt{siitperf} only with flow control (in the Gigabit Ethernet environment) makes the comparison of their results more difficult.

We plant to purchase a NetFPGA device like the one used by the authors of [8] and examine the inter-frame time of the traffic generated by \texttt{siitperf}.

We also plan to test the accuracy of \texttt{siitperf} in a 10GBase-T environment with a Spirent SPT-N4U Tester used out of courtesy for the measurements of [17].

7. Conclusion

We have carefully examined, what kind of factors may distort the measurement results of \texttt{siitperf}, and we set up an error model.
We have compared the results of \texttt{siitperf} used with and without Ethernet flow control in a 10GBase-T environment, and we found that the deviation of the results was always below 1%.

We have calibrated \texttt{siitperf} with a commercial Tester in a Gigabit Ethernet environment, and we found that the deviation of the results was below 1%.

We conclude that it is necessary to calibrate \texttt{siitperf} also in a 10GBase-T environment and we plan to do so.

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